Simplified and Improved Bounds on the VC-Dimension for Elastic Distance Measures

## Frederik Brüning Anne Driemel

University of Bonn

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## VC-Dimension of a Range Space

Range space $\mathcal{R}$ with ground set $X$ :
Set $\mathcal{R}$ s.t. any $r \in \mathcal{R}$ is of the form $r \subseteq X$

## Example:

Balls with ground set $X=\mathbb{R}^{2}$ :

$$
\mathcal{R}=\left\{b(c, \Delta) \mid \Delta \in \mathbb{R}_{+}, c \in \mathbb{R}^{2}\right\}
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where

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b(c, \Delta)=\left\{x \in \mathbb{R}^{2} \mid\|x-c\| \leq \Delta\right\}
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$=$

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Shattering:
$A \subseteq X$ is shattered by $\mathcal{R} \Longleftrightarrow$
$\forall A^{\prime} \subseteq A \exists r \in \mathcal{R}$ s.t. $A^{\prime}=r \cap A$
(all subsets of $A$ can be realized by ranges in $\mathcal{R}$ )
VC-dim $(\mathcal{R})$ :
Maximal size of a shattered subset $A \subseteq X$

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## Motivation: What are VC-Dimension bounds used for?

Sample bounds for computational tasks:

- $\epsilon$-nets, relative-error ( $p, \epsilon$ )-approximations
- test error of classification model

Applications:

- kernel density estimation
- coresets
- clustering
- object recognition


## Range Spaces for Elastic Distance Measures

Range spaces:

$$
\begin{aligned}
d_{\rho} & =\text { distance measure on }\left(\mathbb{R}^{d}\right)^{m} \\
b_{\rho}(c, \Delta) & =\left\{x \in\left(\mathbb{R}^{d}\right)^{m} \mid d_{\rho}(x, c) \leq \Delta\right\} \\
\mathcal{R}_{\rho, k} & =\left\{b_{\rho}(c, \Delta) \mid \Delta \in \mathbb{R}_{+}, c \in\left(\mathbb{R}^{d}\right)^{k}\right\}
\end{aligned}
$$

Distance measures:

- Hausdorff
- (weak) Fréchet
- Dynamic Time Warping

Ground set $\left(\left(\mathbb{R}^{d}\right)^{m}\right) /$ Centers $\left(\left(\mathbb{R}^{d}\right)^{k}\right)$ :

- Polygonal regions
- continuous polygonal curves
- discrete polygonal curves




## Hausdorff Distance

directed Hausdorff distance:
$d_{\vec{H}}(P, Q)=\sup _{p \in P} \inf _{q \in Q}\|p-q\|$
Hausdorff distance:

$$
d_{H}(P, Q)=\max \left\{d_{\vec{H}}(P, Q), d_{\vec{H}}(Q, P)\right\}
$$

for $P, Q \subseteq \mathbb{R}^{d}$


## Results

|  |  | new | old |
| :---: | :---: | :---: | :---: |
| discrete <br> polygonal <br> curves | DTW | $O\left(d k^{2} \log (m)\right)$ | - |
|  | Hausdorff | $O(d k m \log (k))$ |  |
|  | Fréchet | $O(d k \log (k m))$ | $O(d k \log (k m))^{(*)}$ |
| continuous <br> polygonal <br> curves | Hausdorff | $O(d k \log (k m))$ |  |
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Table: Overview of VC-dimension bounds. Results marked with ${ }^{(*)}$ were independently obtained by Cheng and Huang [2024]. The old results were obtained by Driemel, Nusser, Phillips and Psarros [2021].

Lower bound (Driemel, Nusser, Phillips and Psarros [2021]): $\Omega(\max (d k \log (k), \log (d m)))$ for $d \geq 4$ for polygonal curves.

## General Approach

## Approach:

- Split query $d_{H}(P, Q)$ into predicates
- Express predicates as combinations of sign values of polynomials
- Bound VC-dim based on number of cells in arrangement of zero sets of polynomials.


## Definition:

F: Class of polynomials of constant degree from $\mathbb{R}^{d k+1} \times \mathbb{R}^{d m}$ to $\mathbb{R}$
$\mathcal{R}$ is $t$-combination of $\operatorname{sgn}(F): \exists$ boolean function $g, \exists f_{1}, \ldots, f_{t} \in F$ s.t. $\forall r \in \mathcal{R} \exists y$ s.t.
$r=\left\{x \in X \mid g\left(\operatorname{sgn}\left(f_{1}(y, x)\right), \ldots, \operatorname{sgn}\left(f_{t}(y, x)\right)\right)=1\right\}$


Theorem (Anthony and Bartlett 1999):
Suppose $\mathcal{R}$ is a $t$-combination of $\operatorname{sgn}(F)$. Then $\operatorname{VC-dim}(\mathcal{R})$ is in $O(d k \log (t))$. Idea goes back to Goldberg and Jerrum [1993] and independently Ben-David and Lindenbaum [1993].

## Predicates



Predicates $\mathcal{P}_{1}$ and $\mathcal{P}_{2}$ are by Driemel, Nusser, Phillips and Psarros [2021]

## Predicates

## Cases:

$d_{\vec{H}}(P, Q)$ maximized at point $p$ at the boundary of $P$
$d_{\vec{H}}(Q, P)$ maximized at point $q$ in the interior of $Q$


## Predicates:

- $(\mathcal{B})$ (Boundary): True $\Longleftrightarrow d_{\vec{H}}(\partial P, Q) \leq \Delta$.
- (I) (Interior): True if $d_{\vec{H}}(P, Q) \leq \Delta$. False if $d_{\vec{H}}(P, Q)>d_{\vec{H}}(\partial P, Q)$ and $d_{\vec{H}}(P, Q)>\Delta$.

$$
d_{\vec{H}}(P, Q) \leq \Delta \Longleftrightarrow(\mathcal{B}) \text { and }(\mathcal{I}) \text { true }
$$

## Predicate ( $\mathcal{B}$ )


(B) (Boundary): True $\Longleftrightarrow d_{\vec{H}}(\partial P, Q) \leq \Delta$
$d_{\vec{H}}(\partial P, Q) \leq \Delta \Longleftrightarrow d_{\vec{H}}(e, Q) \leq \Delta$ for every edge $e$ of $P$

- Find for each $e$ the part that is outside of $Q$ (here $\overline{s t}$ )
- Find sequence of edges of $Q$ such that these parts are included in their stadiums (here $a, b, c$ )


## Predicate ( $\mathcal{B}$ )

Find part of $e$ that is outside of $Q$ (here $\overline{s t}$ )


- $\mathcal{P}_{3}: p \in Q$ ?
- $\mathcal{P}_{4}: e_{1} \cap e_{2} \neq \emptyset$ ?
- $\mathcal{P}_{5}: e_{1}$ intersects $e_{2}$ before $e_{3}$ ?



## Predicate ( $\mathcal{B}$ )

Find stadiums that include $s$ and $t$ (here $a$ and $c$ )


- $\mathcal{P}_{6} / P_{7}: \exists b$ on $e_{3}$ with $\|b-a\| \leq \Delta$ ?
- $\mathcal{P}_{6}: a$ is first point on $e_{1}$ in $e_{1} \cap e_{2}$
- $\mathcal{P}_{7}: a$ is last point on $e_{1}$ in $e_{1} \cap e_{2}$



## Predicate ( $\mathcal{B}$ )

Find stadiums that include $s$ and $p_{2}$ (here $a$ and $c$ )


- $\mathcal{P}_{1}: \exists q$ on $e_{1}$ with $\|p-q\| \leq \Delta$ ?



## Predicate ( $\mathcal{B}$ )

Find sequence of stadiums that include $\overline{s t}$ (here $a, b, c$ )


- $D_{\Delta, 2}\left(e_{1}, e_{2}\right)$ : intersection of stadiums around $e_{1}, e_{2}$
- $\mathcal{P}_{2}: \ell\left(\overline{p_{j} p_{j+1}}\right) \cap D_{\Delta, 2}\left(e_{1}, e_{2}\right) \neq \emptyset$ ?



## Predicate ( $\mathcal{I}$ )

$(\mathcal{I})$ (Distance realized in interior): We only check vertices of Voronoi diagram of edges of $Q$.


## Predicate (I)

$(\mathcal{I})$ (Distance realized in interior): We only check vertices of Voronoi diagram of edges of $Q$.


c)


We check distances of all Voronoi vertices to all edges of $Q$.
A Voronoi vertex is relevant if it is inside of $P$ and outside of $Q$.


- $\mathcal{P}_{8}: \exists p$ on $e_{4}$ with $\|v-p\| \leq \Delta$ ?
- $\mathcal{P}_{9}: v \in Q$ ?
- $\mathcal{P}_{10}: v \in P$ ?


## Remaining step

## Lemma

For any two polygonal regions $P$ and $Q$ (that may contain holes), given the truth values of all predicates of the type $\mathcal{P}_{1}, \ldots, \mathcal{P}_{10}$ one can determine whether $d_{\vec{H}}(P, Q) \leq \Delta$.

## Remaining step:

Express predicates $\mathcal{P}_{1}, \ldots, \mathcal{P}_{10}$ as combinations of sign values of polynomials.

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